



Algorithms: Design
and Analysis, Part II

Exact Algorithms for NP-Complete Problems

The Vertex Cover Problem

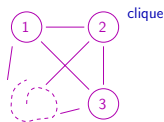
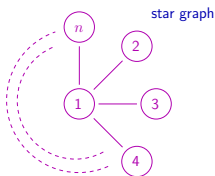
The Vertex Cover Problem

Input: An undirected graph $G = (V, E)$.

Goal: Compute a minimum-cardinality **vertex cover** – a subset $S \subseteq V$ that contains at least one endpoint of each edge of G .

Quiz

Question: What is the minimum size of a vertex cover of a star graph with n vertices and a clique with n vertices respectively?



- A) 1 and $n - 1$
- B) 1 and n
- C) 2 and $n - 1$
- D) $n - 1$ and n

Fact: In general, Vertex Cover is an NP-complete problem.

Strategies for NP-Complete Problems

(1) Identify computationally tractable special cases

- Trees [application of dynamic programming - try it!]
- Bipartite graphs [application of the maximum flow problem]
- When the optimal solution is “small” ($\approx \log n$ or less)

(2) Heuristics (e.g., via suitable greedy algorithms)

(3) Exponential time but better than brute-force search [coming up next]