

Design and Analysis  
of Algorithms I

# Data Structures

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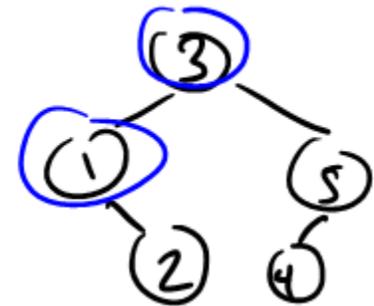
## Red-Black Trees

# Binary Search Tree Structure

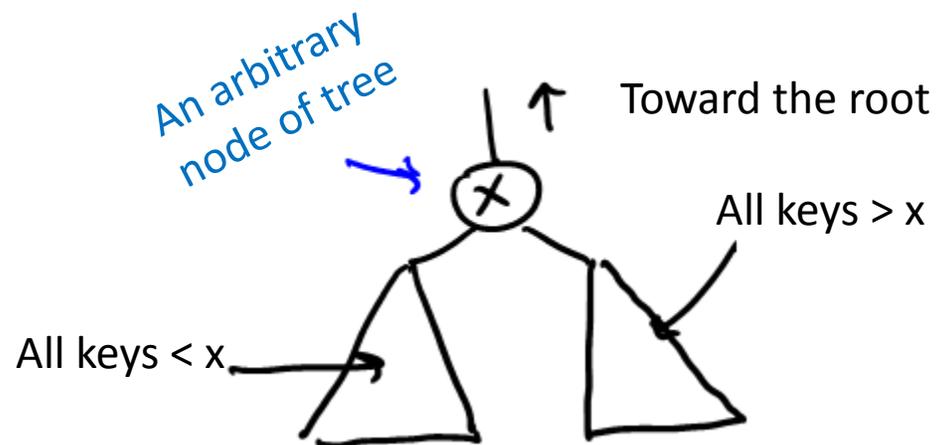
- exactly one node per key
- most basic version :
  - each node has
    - left child pointer
    - right child pointer
    - parent pointer

**SEARCH TREE PROPERTY :**  
( should hold at every node of the search tree )

Root



Leaves



# The Height of a BST

**Note** : many possible trees for a set of keys.

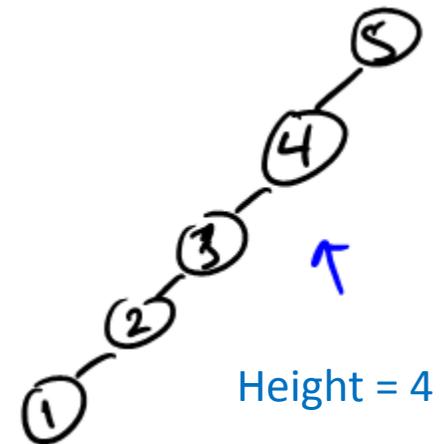
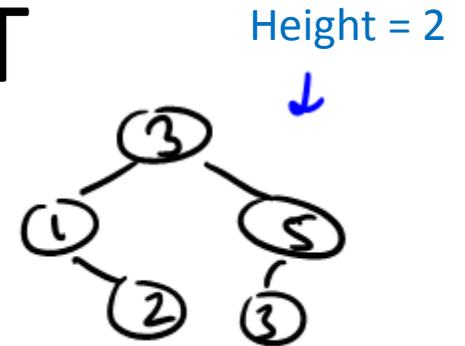
**Note** : height could be anywhere from

$\sim \log_2 n$  to  $\sim n$

Worst case,  
a chain

Best case,  
perfectly  
balanced

(aka depth) longest  
root-leaf path



# Balanced Search Trees

Idea : ensure that height is always  $O(\log(n))$  [best possible]  
 $\Rightarrow$  Search / Insert / Delete / Min / Max / Pred / Succ will then run  
in  $O(\log(n))$  time [n = # of keys in tree]

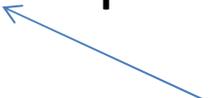
Example : red-black trees [Bayes '72, Guibas-Sedgwick '78]

[ see also AVL trees, splay trees, B trees ]

# Red-Black Invariants

1. Each node red or black
2. Root is black
3. No 2 reds in a row  
[ red node => only black children ]
4. Every root-NULL path has same number of black nodes

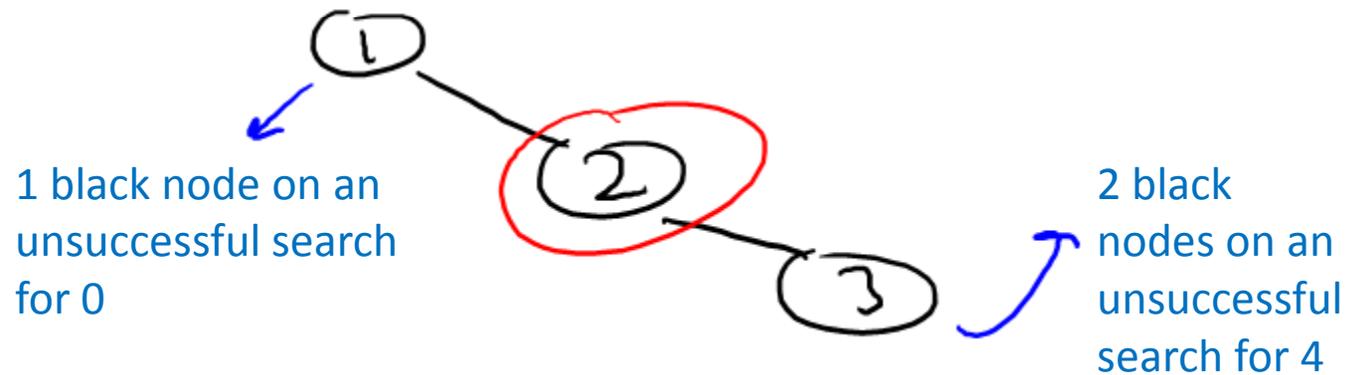
Like in an  
unsuccessful search



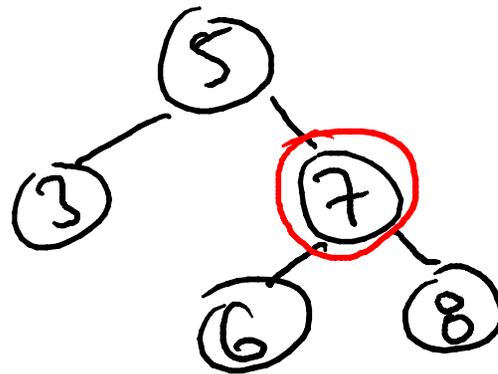
# Example #1

Claim : a chain of length 3 cannot be a red-black tree

Proof



# Example #2

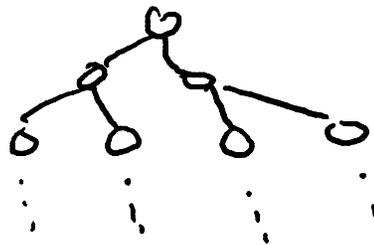


# Height Guarantee

Claim : every red-black tree with  $n$  nodes has height  $\leq 2 \log_2(n + 1)$

Proof : Observation : if every root-NUL path has  $\geq k$  nodes, then tree includes (at the top) a perfectly balanced search tree of depth  $k-1$ .

$\Rightarrow$  Size  $n$  of the tree must be at least  $2^k - 1$



[ $k = 3$ ]

# Height Guarantee (con'd)

Story so far : size  $n \geq 2^k - 1$  , where  $k$  = minimum # of nodes on root – NULL path

$\Rightarrow k \leq \log_2(n + 1)$

Thus : in a red-black tree with  $n$  nodes, there is a root-NULL path with at most  $\log_2(n + 1)$  black nodes.

By 4<sup>th</sup> Invariant : every root-NULL path has  $\leq \log_2(n + 1)$  black nodes

By 3<sup>rd</sup> Invariant : every root-NULL path has  $\leq 2 \log_2(n + 1)$  total nodes.

**Q.E.D.**

Which of the search tree operations have to be re-implemented so that the Red-Black invariants are maintained?

- Search
- Delete
- Insert and Delete
- None of the above